

## 4. TENSOR PRODUCT

Here is an outline of what I did:

- (1) categorical definition
- (2) construction
- (3) list of basic properties
- (4) distributive property
- (5) right exactness
- (6) localization is flat
- (7) extension of scalars
- (8) applications

4.1. **definition.** First I gave the categorical definition and then I gave an explicit construction.

4.1.1. *universal condition.* Tensor product is usually defined by the following universal condition.

**Definition 4.1.** If  $E, F$  are two modules over a commutative ring  $R$ , their *tensor product*  $E \otimes F$  is defined to be the  $R$ -module having the following universal property. First, there exists an  $R$ -bilinear mapping

$$f : E \times F \rightarrow E \otimes F.$$

Second, this mapping is universal in the sense that, for any other  $R$ -module  $M$  and bilinear mapping  $g : E \times F \rightarrow M$ , there exists a unique  $R$ -module homomorphism  $h : E \otimes F \rightarrow M$  making the following diagram commute.

$$\begin{array}{ccc} E \times F & \xrightarrow{f} & E \otimes F \\ & \searrow g & \swarrow \exists! h \\ & & M \end{array}$$

As with all universal conditions, this definition only gives the uniqueness of  $E \otimes F$  up to isomorphism. For the existence we need a construction.

4.1.2. *construction of  $E \otimes F$ .* The mapping  $f : E \times F \rightarrow E \otimes F$  is not onto. However, the image must generate  $E \otimes F$  otherwise we get a contradiction. The elements in the image of  $f$  are denoted

$$f(x, y) = x \otimes y.$$

**Definition 4.2.** The tensor product  $E \otimes F$  is defined to be the  $R$  module which is generated by the symbols  $x \otimes y$  for all  $x \in E, y \in F$  modulo the following conditions

- (1)  $x \otimes -$  is  $R$ -bilinear. I.e.
- (a)  $x \otimes ry = r(x \otimes y)$  for all  $r \in R$
  - (b)  $x \otimes (y + z) = (x \otimes y) + (x \otimes z)$
- (2)  $- \otimes y$  is  $R$ -bilinear. I.e.,
- (a)  $rx \otimes y = r(x \otimes y)$  for all  $r \in R$
  - (b)  $(x + y) \otimes z = (x \otimes z) + (y \otimes z)$

I pointed out that these conditions require  $R$  to be commutative since

$$rs(x \otimes y) = r(sx \otimes y) = sx \otimes ry = s(x \otimes ry) = sr(x \otimes y).$$

**Proposition 4.3.**  $E \otimes F$  as given in the second definition satisfies the universal condition of the first definitions and therefore, the tensor product exists and is unique up to isomorphism.

*Proof.* I said in class that this is obvious. If there is a bilinear mapping  $g : E \times F \rightarrow M$ , the induced mapping  $h : E \otimes F \rightarrow M$  must take the generators  $x \otimes y$  to  $g(x, y)$ . Otherwise the diagram will not commute. Therefore,  $h$  is given on the generators and is thus unique. The only thing we need is to show that  $h$  is a homomorphism. But this is equivalent to showing that the elements of the form

$$rx \otimes y - r(x \otimes y)$$

and elements corresponding to the other three conditions in the second definition go to zero in  $M$ . But this element goes to

$$g(rx, y) - rg(x, y) = 0$$

since  $g$  is  $R$ -bilinear and similarly for the other three elements. So,  $h$  is an  $R$ -module homomorphism and we are done.  $\square$

4.1.3. *functorial properties of tensor product.* The first properties I mentioned were the categorical properties which follow directly from the definition.

**Proposition 4.4.** For a fixed  $R$ -module  $M$ , tensor product with  $M$  is a functor

$$M \otimes - : R\text{-Mod} \rightarrow R\text{-Mod}.$$

What this means is that, given an homomorphism  $f : A \rightarrow B$  there is an  $R$ -module homomorphism

$$1 \otimes f : M \otimes A \rightarrow M \otimes B$$

which satisfies two conditions:

- (1)  $1 \otimes id_A = id_{M \otimes A}$

$$(2) 1 \otimes fg = (1 \otimes f)(1 \otimes g).$$

The definition is  $(1 \otimes f)(x \otimes y) = x \otimes f(y)$ . This gives a homomorphism since the mapping  $M \times A \rightarrow M \otimes B$  given by

$$(x, y) \mapsto x \otimes f(y)$$

is bilinear and therefore induces the desired mapping  $1 \otimes f$ .

More generally, given two homomorphisms  $f : M \rightarrow N, g : A \rightarrow B$  we get a homomorphism

$$f \otimes g : M \otimes A \rightarrow N \otimes B$$

by the formula

$$(f \otimes g)(x \otimes y) = f(x) \otimes g(y).$$

**4.2. exact functors and flat modules.** Flat modules are those for which the functor  $M \otimes -$  is exact. An exact functor is one that takes short exact sequences to short exact sequences. So, first I explained the definitions.

**Definition 4.5.** An *exact sequence* is a sequence of modules and homomorphisms so that the image of each map is equal to the kernel of the next map. A *short exact sequence* is an exact sequence of the following form:

$$0 \rightarrow A \xrightarrow{\alpha} B \xrightarrow{\beta} C \rightarrow 0.$$

In other words,  $\alpha : A \rightarrow B$  is a monomorphism,  $\beta : B \rightarrow C$  is an epimorphism and  $\text{im } \alpha = \ker \beta$  or:  $C \cong B/\alpha A$ .

Sometimes short exact sequences are written:

$$A \twoheadrightarrow B \twoheadrightarrow C.$$

**Definition 4.6.** A functor  $F : R\text{-Mod} \rightarrow R\text{-Mod}$  is called *exact* if it takes short exact sequences to short exact sequences. Thus the short exact sequence above should give the short exact sequence

$$0 \rightarrow FA \xrightarrow{F\alpha} FB \xrightarrow{F\beta} FC \rightarrow 0.$$

**Definition 4.7.** An  $R$ -module  $M$  is called *flat* if  $M \otimes -$  is an exact functor. I.e.,

$$0 \rightarrow M \otimes A \xrightarrow{1 \otimes \alpha} M \otimes B \xrightarrow{1 \otimes \beta} M \otimes C \rightarrow 0$$

is exact for all short exact sequences  $A \twoheadrightarrow B \twoheadrightarrow C$ .

One of the main results (which we will see is actually trivial) is that  $S^{-1}R$  is flat for any multiplicative set  $S$ . I.e., localization is exact.

**4.3. list of properties.** I explained that the exactness of localization was one of the key ideas. However, the explanation required an understanding of the basic properties of tensor product. So, I went back to the beginning with this list.

- (0) (unity)  $R \otimes M \cong M$ .
- (1) (commutative)  $M \otimes N \cong N \otimes M$
- (2) (distributive)  $N \otimes \bigoplus M_i \cong \bigoplus (N \otimes M_i)$
- (3) (associative)  $(A \otimes B) \otimes C \cong A \otimes (B \otimes C)$
- (4) (right exactness)  $M \otimes -$  is right exact, i.e., a short exact sequence  $A \rightarrow B \rightarrow C$  gives an exact sequence

$$M \otimes A \xrightarrow{1 \otimes \alpha} M \otimes B \xrightarrow{1 \otimes \beta} M \otimes C \rightarrow 0$$

- (5) (localization is exact) I.e., we get an exact sequence:

$$0 \rightarrow S^{-1}A \rightarrow S^{-1}B \rightarrow S^{-1}C \rightarrow 0.$$

- (6) (extension of scalars) Given a ring homomorphism  $R \rightarrow S$ , every  $R$ -module  $M$  gives an  $S$  module  $S \otimes_R M$ .

4.3.1. *Grothendieck ring.* I did not prove properties (1) and (3). I said they were obvious. However, I put the first three conditions into a conceptual framework by pointing out that these are the axioms of a ring. The only thing that we don't have is an additive inverse. The algebraic construction is as follows.

First, you take the set of all isomorphism classes of finitely generated  $R$ -modules  $[M]$ . This set has addition and multiplication given by

$$[M] + [N] = [M \oplus N]$$

$$[M][N] = [M \otimes N]$$

Addition and multiplication are associative and commutative and have units:  $[0]$  is the additive unit and  $[R]$  is the multiplicative unit. It just doesn't have additive inverses. So, Grothendieck said to just put in formal inverses:

$$[M] - [N]$$

which are defined like fractions:

$$[M] - [N] = [A] - [B]$$

if there exists another module  $C$  so that

$$M \oplus B \oplus C \cong N \oplus A \oplus C.$$

This gives a ring whose name is  $G(R)$ . The notation  $K_0(R)$  is for the ring of formal differences of f.g. projective  $R$ -modules.

4.3.2.  $R \otimes M \cong M$ . After using this formula many times in the lecture, I decided I should prove it. I put the proof at the beginning in the notes where it belongs.

**Theorem 4.8.**  $R \otimes M \cong M$  for any  $R$ -module  $M$ .

*Proof.* Since the mapping

$$R \times M \rightarrow M$$

given by  $(r, x) \mapsto rx$  is bilinear it induces a mapping

$$\mu : R \otimes M \rightarrow M$$

so that  $\mu(r \otimes x) = rx$ . The inverse mapping  $\phi : M \rightarrow R \otimes M$  is given by  $\phi(x) = 1 \otimes x$ . We carefully checked that these are inverse to each other:

$$\phi\mu(r \otimes x) = \phi(rx) = 1 \otimes rx = r(1 \otimes x) = r \otimes x$$

$$\mu\phi(x) = \mu(1 \otimes x) = 1x = x.$$

So, these maps are both isomorphisms of  $R$ -modules.  $\square$

4.3.3. *distributive property.* I gave a category theory proof of the distributivity of tensor product over direct sum. First I pointed out that the following formal characterization of direct sum.

**Lemma 4.9.**  $M$  is the direct sum of modules  $M_1, \dots, M_n$  if and only if there are inclusion maps  $s_i : M_i \rightarrow M$  and projection maps  $p_i : M \rightarrow M_i$  so that

- (1)  $p_j \circ s_i = \delta_{ij}$ , i.e., equal to the identity mapping on  $M_i$  if  $i = j$  and equal to 0 if  $i \neq j$ .
- (2)  $\sum s_i \circ p_i = id_M$ .

I drew the following diagrams to illustrate these equations.

$$\begin{array}{ccc}
 M_i & \xrightarrow{\delta_{ij}} & M_j \\
 & \searrow s_i & \nearrow p_j \\
 & & M
 \end{array}
 \quad p_j \circ s_i = \delta_{ij}$$
  

$$\begin{array}{ccc}
 M & \xrightarrow{\quad} & M \\
 & \searrow p_i & \nearrow s_i \\
 & & M
 \end{array}
 \quad \sum_{i=1}^n s_i \circ p_i = id_M$$

This lemma was proved in any preadditive category in Part B, Theorem 7.4.

**Theorem 4.10.** *If  $M \cong \bigoplus_{i=1}^n M_i$  then*

$$N \otimes M = N \otimes \bigoplus_{i=1}^n M_i \cong \bigoplus_{i=1}^n (N \otimes M_i).$$

*Proof.* Consider the homomorphisms:

$$N \otimes M_i \xrightarrow{1 \otimes s_i} N \otimes M \xrightarrow{1 \otimes p_j} N \otimes M_j$$

- a)  $(1 \otimes p_j)(1 \otimes s_i) = 1 \otimes p_j s_i = 1 \otimes \delta_{ij} = \delta_{ij}(1 \otimes 1)$ .  
 b)  $\sum (1 \otimes s_i)(1 \otimes p_i) = 1 \otimes \sum s_i p_i = 1 \otimes 1 = id_{N \otimes M}$ .

These conditions imply that  $N \otimes M \cong \bigoplus N \otimes M_i$  by the above lemma.  $\square$

*Remark 4.11.* This proof works in any preadditive category to show that any linear functor distributes over direct sum.

**4.4. right exactness of tensor product.** I didn't prove the right exactness of tensor product this first time because the elementary proof is messy and not very instructive. I just explained that this is a special case of a much more general principle that: "All linear left adjoint functors are right exact." I will explain this later. The statement of the theorem is the following.

**Theorem 4.12.** *Tensor product with any  $R$ -module  $M$  sends any exact sequence of  $R$ -modules of the form:*

$$A \xrightarrow{\alpha} B \xrightarrow{\beta} C \rightarrow 0$$

*to another exact sequence of the same form:*

$$M \otimes A \rightarrow M \otimes B \rightarrow M \otimes C \rightarrow 0.$$

This statement appears stronger than the original statement since the hypothesis is weaker. But I explained that the first statement implies this second version. Suppose that we know that  $M \otimes -$  sends short exact sequences to right exact sequences as above. Then how can we conclude that it sends the more general right exact sequences  $A \rightarrow B \rightarrow C \rightarrow 0$  to right exact sequences?

The first statement implies that  $M \otimes -$  takes epimorphisms to epimorphisms. (In fact this is obvious since the generators  $x \otimes y \in M \otimes C$  come from generators  $x \otimes \tilde{y} \in M \otimes B$ .) Therefore  $M \otimes A$  maps onto  $M \otimes \alpha(A)$ . If we assume the weaker condition that the functor  $M \otimes -$

takes short exact sequences to right exact sequences, then it will take the short exact sequence

$$0 \rightarrow \alpha(A) \hookrightarrow B \xrightarrow{\beta} C \rightarrow 0$$

to an exact sequence

$$M \otimes \alpha(A) \rightarrow M \otimes B \xrightarrow{1 \otimes \beta} M \otimes C \rightarrow 0$$

This says that  $M \otimes \alpha(A)$  maps onto the kernel of  $1 \otimes \beta$ . But  $M \otimes A$  maps onto  $M \otimes \alpha(A)$ . So,  $M \otimes A$  also maps onto  $\ker(1 \otimes \beta)$ . So, we get an exact sequence

$$M \otimes A \xrightarrow{1 \otimes \alpha} M \otimes B \xrightarrow{1 \otimes \beta} M \otimes C \rightarrow 0.$$

Here is an example of how this is used.

**Corollary 4.13.** *Suppose that  $I \subset R$  is an ideal. Then*

$$R/I \otimes M \cong M/IM$$

where  $IM$  is the submodule of  $M$  generated by all products of the form  $ax$  where  $a \in I$  and  $x \in M$ . In particular, when  $I = (p)$  is principal, we have

$$R/(p) \otimes M \cong M/pM$$

where  $pM = \{px \mid x \in M\}$ .

*Proof.* Suppose that  $I$  is generated by elements  $a_i$ . Then we have an epimorphism of  $R$  modules

$$\bigoplus_i R \twoheadrightarrow I$$

sending  $(r_i) \in \bigoplus_i R$  to  $\sum r_i a_i \in I$ . This gives an exact sequence

$$\bigoplus_i R \xrightarrow{\alpha} R \rightarrow R/I \rightarrow 0.$$

Tensor with  $M$  to give

$$\bigoplus_i R \otimes M \xrightarrow{\alpha \otimes 1} R \otimes M \rightarrow R/I \otimes M \rightarrow 0.$$

Using the isomorphisms  $\mu : R \otimes M \cong M$  and  $\phi : M \cong R \otimes M$  we get an exact sequence

$$\bigoplus_i M \xrightarrow{\mu(\alpha \otimes 1)\phi} M \rightarrow R/I \otimes M \rightarrow 0$$

where  $\mu(\alpha \otimes 1)\phi$  sends  $(x_i) \in \bigoplus_i M$  to  $\sum a_i x_i \in M$ . The image is equal to  $IM$  by definition. So,  $R/I \otimes M \cong M/IM$  as claimed.  $\square$

For finitely generated modules over a PID we can now compute the tensor product:

$$M \otimes \left( R^n \oplus \bigoplus R/(p_i^{n_i}) \right) \cong M^n \oplus \bigoplus M/p_i^{n_i} M.$$

**4.5. localization is exact.** Recall that a *multiplicative set* is a subset  $S \subseteq R$  which is closed under multiplication, contains 1 and does not contain 0. The *localization*  $S^{-1}R$  was defined to be the ring of all fractions  $r/s$  where  $r \in R$  and  $s \in S$  modulo the equivalence relation

$$\frac{r}{s} \sim \frac{r'}{s'}$$

if there is an element  $t \in S$  so that  $rs't = r'st$ . This ring is also an  $R$ -module since we have an action of  $R$  given by

$$r \cdot \frac{x}{s} = \frac{rx}{s}.$$

**Proposition 4.14.** *For any  $R$ -module  $M$  let  $S^{-1}M$  be the set of equivalence classes of fractions  $x/s$  where  $x \in M, s \in S$  modulo the equivalence relation  $x/s \sim y/s'$  if there is a  $t \in S$  so that  $ts'x = tsy$ . Then  $S^{-1}M$  is an  $R$ -module with action of  $R$  given by  $r(x/s) = rx/s$  and*

$$S^{-1}M \cong S^{-1}R \otimes M.$$

*Proof.* There is an obvious map  $S^{-1}R \otimes M \rightarrow S^{-1}M$  sending  $r/s \otimes x$  to  $rx/s$ . The inverse map sends  $x/s$  to  $1/s \otimes x$ . To show that this is well-defined, take an equivalent element  $tx/ts$ . This goes to

$$\frac{1}{ts} \otimes tx = t \left( \frac{1}{ts} \otimes x \right) = \frac{t}{ts} \otimes x = \frac{1}{s} \otimes x.$$

The rest of the proof is straightforward. □

**Theorem 4.15.**  *$S^{-1}R$  is a flat  $R$ -module. Equivalently, every short exact sequence of  $R$ -modules  $A \rightarrow B \rightarrow C$  induces an exact sequence*

$$0 \rightarrow S^{-1}A \rightarrow S^{-1}B \rightarrow S^{-1}C \rightarrow 0.$$

*Proof.* Since tensor product is right exact, it suffices to show that  $S^{-1}A \rightarrow S^{-1}B$  is a monomorphism. This is easy. We can assume that  $A \subseteq B$  and suppose that  $a \in A$  and  $s \in S$  so that the element  $a/s \in S^{-1}A$  goes to zero in  $S^{-1}B$ . This means

$$\frac{a}{s} \sim \frac{0}{s}$$

in  $S^{-1}B$ . By definition this is equivalent to saying that there exists  $t \in S$  so that  $tsa = 0$ . But this same equation implies that  $a/s = 0/1$  in  $S^{-1}A$ . So, we are done.  $\square$

The ring  $S^{-1}R$  acts on the module  $S^{-1}M$  in the obvious way:

$$\frac{r}{s} \frac{x}{t} = \frac{rx}{st}.$$

This makes  $S^{-1}M$  into a module over  $S^{-1}R$ . This is an example of “extension of scalars.”

**4.6. extension of scalars.** We had a concept before called “restriction of scalars.” That was when we had a subring  $S$  of  $R$  or, more generally, a ring homomorphism  $\phi : S \rightarrow R$  and we got an induced map

$$\phi^* : R\text{-Mod} \rightarrow S\text{-Mod}$$

which sent an  $R$ -module  $M$  to the same thing with the action of  $S$  given by  $s \cdot x = \phi(s)x$ . I.e., we restricted the action of the ring to  $S$ .

“Extension of scalars” goes the other way.

**Proposition 4.16.** *Given a ring homomorphism  $\phi : R \rightarrow S$  and an  $R$ -module  $M$ ,  $S \otimes_R M$  is an  $S$ -module with action of  $S$  given by*

$$s(t \otimes x) = st \otimes x.$$

The module is sometimes written as  $S \otimes_\phi M$  because the  $R$ -module structure is given by

$$r(s \otimes x) = (\phi(r)s) \otimes x = s \otimes rx.$$

This is the  $R$ -module structure induced from the  $S$ -module structure by restriction of scalars.

*Proof.* Multiplication by elements of  $S$  gives an  $R$ -linear map  $S \rightarrow S$  and therefore gives an  $R$ -linear map  $S \otimes M \rightarrow S \otimes M$  by naturality of tensor product. This gives a sequence of ring homomorphisms

$$S \rightarrow \text{End}_R(S) \rightarrow \text{End}_R(S \otimes M)$$

which defines the  $S$ -module structure on  $S \otimes M$ .  $\square$

One special case of this is when  $R$  is a domain and  $F = Q(R)$  is the field of fractions.

**Definition 4.17.** Suppose that  $M$  is a module over a domain  $R$ . Then the *rank* of  $M$  is defined to be the dimension of  $Q(R) \otimes M$  as a vector space over the field  $Q(R)$ .

$$r(M) = \dim_{Q(R)} Q(R) \otimes M.$$

**Theorem 4.18.** *For a f.g. module  $M$  over a PID  $R$ , if*

$$M \cong R^r \oplus \bigoplus R/(p_i^{n_i}),$$

*the number  $r$  is equal to the rank of  $M$  and is therefore uniquely determined.*

*Proof.* This is a calculation using the fact that

$$R/(a) \otimes Q(R) \cong Q(R)/aQ(R) = 0$$

since  $aQ(R) = Q(R)$  for  $a \neq 0$ :

$$Q(R) \otimes M = Q(R) \otimes R^r \oplus \bigoplus Q(R)/p_i^{n_i}Q(R) \cong Q(R)^r.$$

□

It still remains to show that the numbers  $p_i^{n_i}$  are uniquely determined.